

# One-time pad

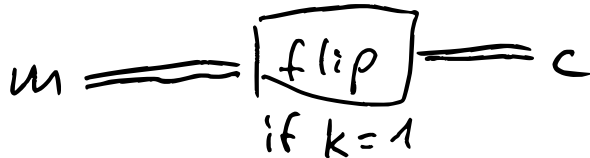
Classically:

Alice wants to send bit  $m$  to Bob with "perfect secrecy"

Secret key:  $k \in \{0, 1\}$

Message  $m \in \{0, 1\}$

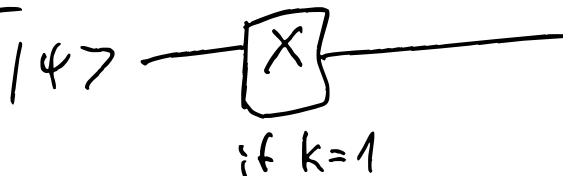
Alice sends:  $c = m \oplus k$



Quantumly:

Message: qubit

1. try



Secure?

$$|0\rangle \rightsquigarrow \left\{ \begin{array}{l} \text{if } k=0 \\ |0\rangle @ \frac{1}{\sqrt{2}}, \end{array} \quad \begin{array}{l} \text{if } k=1 \\ |1\rangle @ \frac{1}{\sqrt{2}} \end{array} \right\}$$

$$|1\rangle \rightsquigarrow \left\{ \begin{array}{l} |1\rangle @ \frac{1}{\sqrt{2}}, \\ |0\rangle @ \frac{1}{\sqrt{2}} \end{array} \right\}$$

$$|+\rangle \rightsquigarrow \left\{ \begin{array}{l} |+\rangle @ \frac{1}{\sqrt{2}}, \\ |+\rangle @ \frac{1}{\sqrt{2}} \end{array} \right\} \\ = \left\{ |+\rangle @ 1 \right\}$$

$$|-\rangle \rightsquigarrow \left\{ \begin{array}{l} |-\rangle @ \frac{1}{\sqrt{2}}, \\ |-\rangle @ \frac{1}{\sqrt{2}} \end{array} \right\} \\ \approx \left\{ |-\rangle @ 1 \right\}$$

$\Rightarrow$  Adv can perfectly distinguish  
the enc of  $|+\rangle$  and  $|-\rangle$ .

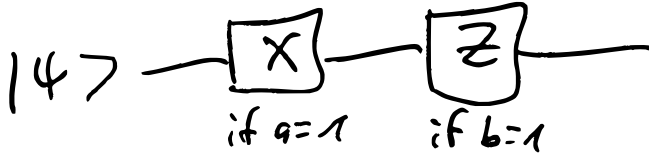
$\Rightarrow$  insecure

$$\begin{array}{l} |+\rangle = \frac{1}{\sqrt{2}}|0\rangle + \frac{1}{\sqrt{2}}|1\rangle \\ \xrightarrow{X} \frac{1}{\sqrt{2}}|1\rangle + \frac{1}{\sqrt{2}}|0\rangle \\ = |+\rangle \end{array}$$

# Quantum OTP

Message:  $|4\rangle \in \mathbb{C}^2$

Key:  $a, b \in \{0, 1\}$



Secure?

Def: For any  $|4\rangle, |4'\rangle$ ,  
the output of this circuit  
is phys. ind.

Thm: Secure

Proof:

Supp  $|4\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$ ,  $|\alpha|^2 + |\beta|^2 = 1$

$E = \left\{ |4\rangle \otimes \frac{1}{4}, X|4\rangle \otimes \frac{1}{4}, Z|4\rangle \otimes \frac{1}{4}, \right.$   
 $\left. ZX|4\rangle \otimes \frac{1}{4} \right\}$

↑  
distr.  
of output

Caveat: This def.  
does not per se  
guarantee anything  
if msg is entangled  
with something else

$$E = \left\{ |4\rangle\langle\frac{1}{4}|, X|4\rangle\langle\frac{1}{4}|, Z|4\rangle\langle\frac{1}{4}|, ZX|4\rangle\langle\frac{1}{4}| \right\}$$

$$S = \frac{1}{4} \left( |4\rangle\langle 4| + X|4\rangle\langle 4|X + Z|4\rangle\langle 4|Z + ZX|4\rangle\langle 4|ZX \right)$$

$$= \frac{1}{4} (A + ZAZ) \text{ with}$$

$$A := |4\rangle\langle 4| + X|4\rangle\langle 4|X$$

$$|4\rangle\langle 4| = \begin{pmatrix} \alpha \\ \beta \end{pmatrix} (\bar{\alpha} \ \bar{\beta}) = \begin{pmatrix} |\alpha|^2 & \alpha\bar{\beta} \\ \bar{\alpha}\beta & |\beta|^2 \end{pmatrix}$$

$$X|4\rangle\langle 4|X = \begin{pmatrix} |\beta|^2 & \bar{\alpha}\beta \\ \alpha\bar{\beta} & |\alpha|^2 \end{pmatrix}$$

$$A = \begin{pmatrix} 1 & \alpha\bar{\beta} + \bar{\alpha}\beta \\ \alpha\bar{\beta} + \bar{\alpha}\beta & 1 \end{pmatrix} \quad ZAZ = \begin{pmatrix} 1 & -\alpha\bar{\beta} - \bar{\alpha}\beta \\ -\alpha\bar{\beta} - \bar{\alpha}\beta & 1 \end{pmatrix}$$

$$A + ZAZ = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix} \quad S = \frac{1}{2} I$$

$$\begin{array}{l} X = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, X^\dagger = X \\ Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}, Z^\dagger = Z \end{array}$$

Output for  $|\psi'\rangle$ : also  $\rho = \frac{1}{2}I$

$\Rightarrow$  Enc  $|\psi\rangle, |\psi'\rangle$

are phys. indist.

$\Rightarrow$  QOTP secure

(perfect secrecy)  $\square$

Thm/Def:  $\rho$  is a density op. iff

-  $\rho$  Hermitian ( $\rho^\dagger = \rho$ )

-  $\rho \geq 0$  (all eigenvalues  $\geq 0$ )

-  $\text{tr } \rho = 1$

$\hookrightarrow$  If we have distrib with total prob  $\leq 1$ , this becomes " $\text{tr } \rho \leq 1$ "

# Throwing away stuff (partial base)

## Motivation 1:

Have  $f: \{0,1\}^n \rightarrow \{0,1\}^m$

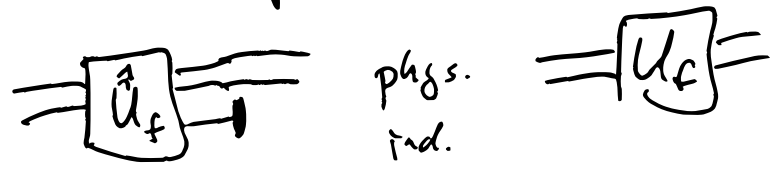
Want:  $U_f: |x, y\rangle \rightarrow |x, y \oplus f(x)\rangle$

Attempt 1:  $\hat{U}_f: |x, y\rangle \rightarrow |x, y \oplus f(x), \text{stuff}\rangle$

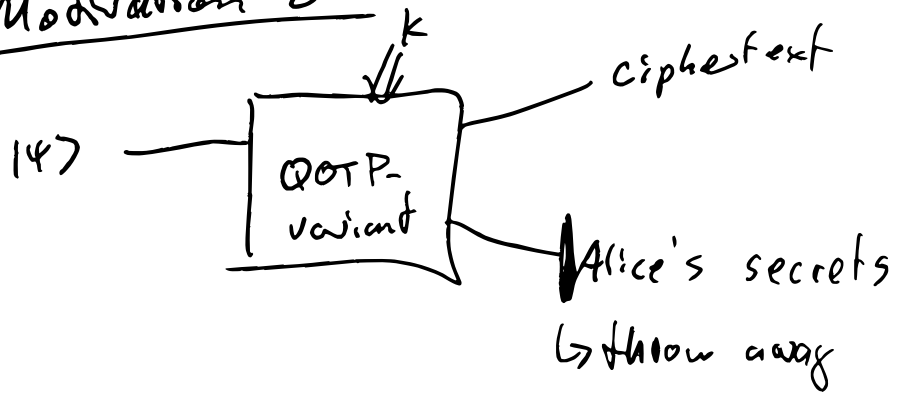
Attempt 2:

$\hat{U}_f: |x, y\rangle \rightarrow |x, y \oplus f(x), 0\rangle$

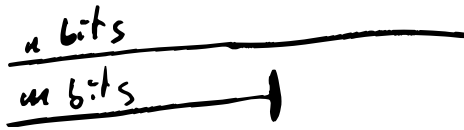
Want to prove:



## Motivation 2:

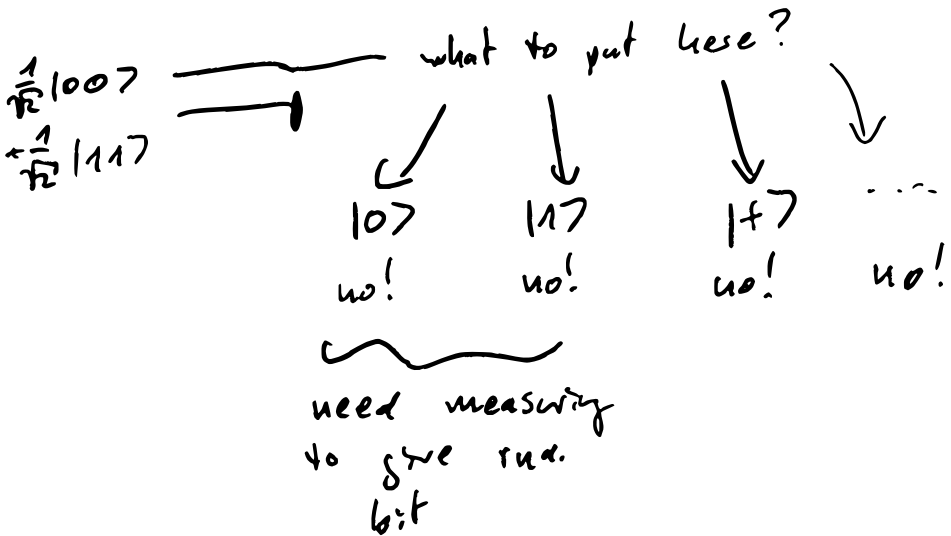


How to mathematically describe:



Say we describe it as a function  $\mathbb{C}^{2^n \cdot 2^m} \rightarrow \mathbb{C}^{2^n}$ .

This does not work!



$\Rightarrow$  Need density op. to describe  $\rightarrow$

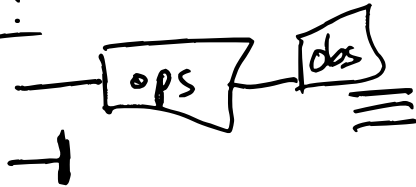
Want:

- some op that maps density ops on  $AB$  ( $\mathbb{C}^{NM \times NM}$ )

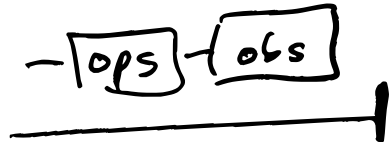
→ density ops on  $A$  ( $\mathbb{C}^{N \times M}$ )

- model "throwing away"

i.e.:



≡



i.e.:



≡  
↑  
same  
density  
op

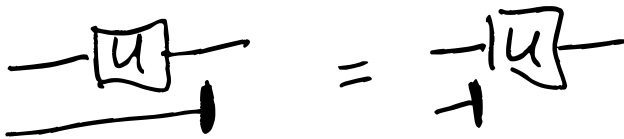




Answer:  $\text{tr}_B \rho$  defined by:

- $\rho \otimes \tau \xrightarrow{\text{tr}_B} \rho \cdot \text{tr} \tau$
  - $\text{tr}_B$  linear
- 

$\text{tr}_B$  "models throwing away" ..



$$\text{i.e.: } \frac{\text{tr}_B (U \otimes I) \rho (U^\dagger \otimes I)}{\text{tr}_B} = U \text{tr}_B U^\dagger$$

and similar for measurements etc.